

# On Faster Convergence of the Bisection Method for Certain Triangles

By Martin Stynes

**Abstract.** Let  $\Delta ABC$  be a triangle with vertices  $A, B$  and  $C$ . It is "bisected" as follows: choose a/the longest side (say  $AB$ ) of  $\Delta ABC$ , let  $D$  be the midpoint of  $AB$ , then replace  $\Delta ABC$  by two triangles,  $\Delta ADC$  and  $\Delta DBC$ .

Let  $\Delta_{01}$  be a given triangle. Bisect  $\Delta_{01}$  into two triangles  $\Delta_{11}, \Delta_{12}$ . Next, bisect each  $\Delta_{1i}, i = 1, 2$ , forming four new triangles  $\Delta_{2i}, i = 1, 2, 3, 4$ . Continue thus, forming an infinite sequence  $T_j, j = 0, 1, 2, \dots$ , of sets of triangles, where  $T_j = \{\Delta_{ji}: 1 \leq i \leq 2^j\}$ . It is known that the mesh of  $T_j$  tends to zero as  $j \rightarrow \infty$ . It is shown here that if  $\Delta_{01}$  satisfies any of four certain properties, the rate of convergence of the mesh to zero is much faster than that predicted by the general case.

**1. Introduction.** Let  $\Delta ABC$  be a triangle with vertices  $A, B$  and  $C$ . We define the procedure for "bisecting"  $\Delta ABC$  as follows: choose a/the longest side (say  $AB$ ) of  $\Delta ABC$ , let  $D$  be the midpoint of  $AB$ , then divide  $\Delta ABC$  into the two triangles  $\Delta ADC$  and  $\Delta DBC$ .

Let  $\Delta_{01}$  be a given triangle. Bisect  $\Delta_{01}$  into two triangles  $\Delta_{11}$  and  $\Delta_{12}$ . Next, bisect each  $\Delta_{1i}, i = 1, 2$ , forming four new triangles  $\Delta_{2i}, i = 1, 2, 3, 4$ . Set  $T_j = \{\Delta_{ji}: 1 \leq i \leq 2^j\}, j = 0, 1, 2, \dots$ , so  $T_j$  is a set of  $2^j$  triangles. Define  $m_j$ , the mesh of  $T_j$ , to be the length of the longest side among the sides of the triangles in  $T_j$ . Clearly  $0 < m_{j+1} \leq m_j$  for all  $j \geq 0$ . It is shown implicitly in [3] that in fact  $m_j \rightarrow 0$  as  $j \rightarrow \infty$ . Thus, this bisection method is useful in finite element methods for approximating solutions of differential equations (see e.g. [1]). A modification of such a bisection method can be used in computing the topological degree of a mapping from  $R^3$  to  $R^3$  ([4], [5]).

In [2] an explicit bound is obtained for the rate of convergence of  $m_j$ :  $m_j \leq (\sqrt{3}/2)^{[j/2]} m_0$ , where  $[x]$  denotes the integer part of  $x$ . In [2] it is also mentioned that computer experiments indicate that in many cases this bound is unrealistically high; this prompted the present results. We show that if  $\Delta_{01}$  belongs to any one of four sets of equivalence classes of triangles, then we have the substantially improved bounds of Corollaries 1 and 2 below. Much of the notation used is taken from [3].

## 2. Results.

**Definition.** Given three positive numbers  $\rho, \sigma, \tau$  such that  $\rho + \sigma + \tau = \pi$ , define  $(\rho, \sigma, \tau)$  to be the set of all triangles whose interior angles are  $\rho, \sigma, \tau$ .

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With this notation we divide the set of all triangles into similarity classes. Note that  $(\rho, \sigma, \tau) = (\sigma, \rho, \tau)$  etc.

*Definition.* Given a triangle  $\Delta$ , define  $d(\Delta)$ , the diameter of  $\Delta$ , to be the length of the longest side of  $\Delta$ .

*Definition.* Given a similarity class  $(\rho, \sigma, \tau)$  choose any triangle  $\Delta \in (\rho, \sigma, \tau)$  and join the midpoint of its longest side to the opposite vertex. The new side ratio  $r$  of  $(\rho, \sigma, \tau)$  is defined to be the length of this new side divided by  $d(\Delta)$ .

*Remark.* The new side ratio is well-defined since it does not depend on the particular  $\Delta$  chosen in  $(\rho, \sigma, \tau)$ . By Lemma 5.2(i) of [4] we have  $0 < r \leq \sqrt{3}/2$  always.

*Definition.* Using the notation of the introduction, an iteration of the bisection method applied to  $\Delta_{00}$  is defined to be the progression from  $T_j$  to  $T_{j+1}$  for any  $j \geq 0$ . A cycle of the bisection method is defined to be two successive iterations, i.e., the progression from  $T_j$  to  $T_{j+2}$  for any  $j \geq 0$ .

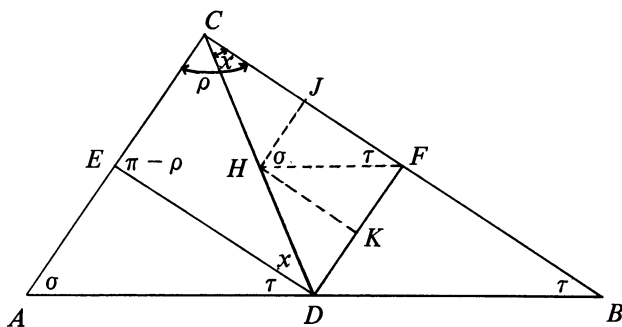


FIGURE 1

In Figure 1,  $\Delta ABC \in (\rho, \sigma, \tau)$  and we have taken  $\tau \leq \sigma \leq \rho$ . Thus,  $AC \leq BC \leq AB$  (where ‘AC’ denotes ‘length of AC’ etc.).  $D, E, F$  are the midpoints of  $AB, AC, BC$ , respectively. Under the additional hypothesis that  $x + \tau \geq \max\{\sigma, \rho - x\}$  and  $\pi - \rho \geq \rho - x$ , bisections will take place exactly as in Figure 1 (see [3]).

*Notation.* To indicate that  $\Delta_1 \in (\alpha_1, \beta_1, \gamma_1)$  yields  $\Delta_2 \in (\alpha_2, \beta_2, \gamma_2)$  and  $\Delta_3 \in (\alpha_3, \beta_3, \gamma_3)$  when bisected, we shall draw a pair of arrows emanating from the triple  $(\alpha_1, \beta_1, \gamma_1)$  with one entering each of the triples  $(\alpha_2, \beta_2, \gamma_2)$  and  $(\alpha_3, \beta_3, \gamma_3)$ .

We quote the following result from [3]; it uses the notation of Figure 1.

LEMMA 1 [3, LEMMA 4]. *If  $\tau \leq \sigma \leq \rho, x + \tau \geq \max\{\sigma, \rho - x\}$ , and  $\pi - \rho \geq \rho - x$ , then we must have*

$$\begin{array}{ccc}
 (\rho, \sigma, \tau) & \begin{array}{c} \longleftarrow \\ \longrightarrow \end{array} & (x, \tau, \rho + \sigma - x) \\
 \downarrow \uparrow & & \downarrow \uparrow \\
 (\rho - x, \sigma, x + \tau) & \begin{array}{c} \longleftarrow \\ \longrightarrow \end{array} & (x, \rho - x, \pi - \rho)
 \end{array}$$

*Proof.* See [3].

Thus, after one cycle of the bisection method applied to  $\Delta ABC \in (\rho, \sigma, \tau)$ , we have two triangles from  $(\rho, \sigma, \tau)$  ( $\Delta ADE, \Delta DBF$  above) and two triangles from  $(x, \rho - x, \pi - \rho)$  ( $\Delta CED, \Delta CFD$  above). Now  $d(ADE) = d(DBF) = d(ABC)/2$ . Also,

$d(CED) = d(CFD) = CD$  (since  $\pi - \rho \geq \rho - x$  by hypothesis, and  $\rho - x \geq x$  from [3]), i.e.,  $d(CED) = d(CFD) = rd(ABC)$ , where  $r$  is the new side ratio of  $(\rho, \sigma, \tau)$ .

Similarly, after one cycle of the bisection method applied to  $\Delta CFD \in (x, \rho - x, \pi - \rho)$ , we obtain  $\Delta HJF, \Delta HKF \in (\rho, \sigma, \tau)$  and  $\Delta CJH, \Delta HKD \in (x, \rho - x, \pi - \rho)$ . Here

$$d(CJH) = d(HKD) = d(CFD)/2,$$

and

$$d(HJF) = d(HKF) = HF = AB/4 = CD/4r = d(CFD)/4r.$$

**THEOREM 1.** Assume that  $\tau \leq \sigma \leq \rho, x + \tau \geq \max\{\sigma, \rho - x\}$ , and  $\pi - \rho \geq \rho - x$ . Then for  $n \geq 1$ , after  $n$  cycles of the bisection method applied to

- (i)  $\Delta \in (\rho, \sigma, \tau)$ , we have  $2^{2n-1}$  triangles in  $(\rho, \sigma, \tau)$  each with diameter  $d(\Delta)/2^n$  and  $2^{2n-1}$  triangles in  $(x, \rho - x, \pi - \rho)$  each with diameter  $d(\Delta)r/2^{n-1}$ ;
- (ii)  $\Delta' \in (x, \rho - x, \pi - \rho)$ , we have  $2^{2n-1}$  triangles in  $(x, \rho - x, \pi - \rho)$  each with diameter  $d(\Delta')/2^n$  and  $2^{2n-1}$  triangles in  $(\rho, \sigma, \tau)$  each with diameter  $d(\Delta')/2^{n+1}r$ .

Here  $r$  is the new side ratio of  $(\rho, \sigma, \tau)$ .

*Proof.* We use induction on  $n$ . The case  $n = 1$  is proven in the remarks following Lemma 1.

Fix  $k > 1$ . Assume that the theorem is true for  $1 \leq n < k$ . We prove it true for  $n = k$ .

First, part (i). After one cycle of the bisection method applied to  $\Delta$ , we have  $\Delta_1, \Delta_2 \in (\rho, \sigma, \tau)$  with  $d(\Delta_1) = d(\Delta_2) = d(\Delta)/2$ , and  $\Delta_3, \Delta_4 \in (x, \rho - x, \pi - \rho)$  with  $d(\Delta_3) = d(\Delta_4) = rd(\Delta)$ . Applying a further  $k - 1$  cycles to each of these four triangles, we obtain from  $\Delta_1$  and  $\Delta_2$  by the inductive hypothesis  $2^{2k-2}$  triangles in  $(\rho, \sigma, \tau)$  each with diameter  $d(\Delta_1)/2^{k-1} = d(\Delta)/2^k$ , and  $2^{2k-2}$  triangles in  $(x, \rho - x, \pi - \rho)$  each with diameter  $rd(\Delta_1)/2^{k-2} = rd(\Delta)/2^{k-1}$ ; from  $\Delta_3$  to  $\Delta_4$  we get  $2^{2k-2}$  triangles in  $(x, \rho - x, \pi - \rho)$  each with diameter  $d(\Delta_3)/2^{k-1} = rd(\Delta)/2^{k-1}$  and  $2^{2k-2}$  triangles in  $(\rho, \sigma, \tau)$  each with diameter  $d(\Delta_3)/r2^k = d(\Delta)/2^k$ . Adding totals of identical triangles shows that (i) holds for  $n = k$ .

By an analogous argument (ii) holds for  $n = k$ . This completes the proof.

**COROLLARY 1.** Suppose  $\tau \leq \sigma \leq \rho, x + \tau \geq \max\{\sigma, \rho - x\}$ , and  $\pi - \rho \geq \rho - x$ . Then in the notation of the introduction

- (i) If  $\Delta_{01} \in (\rho, \sigma, \tau)$ , then  $m_j \leq \max\{r, \frac{1}{2}\}(\frac{1}{2})^{\lfloor j/2 \rfloor - 1} d(\Delta_{01})$  for  $j \geq 1$ , with equality for even  $j$ ;
- (ii) if  $\Delta_{01} \in (x, \rho - x, \pi - \rho)$ , then  $m_j \leq \max\{1/2r, 1\}(\frac{1}{2})^{\lfloor j/2 \rfloor} d(\Delta_{01})$  for  $j \geq 1$ , with equality for even  $j$ .

*Proof.* Immediate from Theorem 1.

*Remark.* In practice the conditions of Theorem 1 are more easily checked if expressed in terms of the lengths of sides of triangles. Using the notation of Figure 1,

$$\tau \leq \sigma \leq \rho \text{ is equivalent to } AC \leq BC \leq AB,$$

$$x + \tau \geq \max\{\sigma, \rho - x\} \text{ is equivalent to } AC \geq \max\{AB/2, CD\},$$

$$\pi - \rho \geq \rho - x \text{ is equivalent to } CD \geq BC/2.$$

Thus, knowing the lengths of  $AC, BC, AB$  and  $CD$  one can immediately decide whether

or not  $\Delta ABC \in (\rho, \sigma, \tau)$  satisfies the conditions of Theorem 1. Note that these inequalities and [4, Lemma 5.2(i)] give  $1/4 \leq r \leq \sqrt{3}/2$ .

Given a triangle such as  $\Delta CFD$  with  $CD \geq CF \geq DF$ , to decide whether or not  $\Delta CFD \in (x, \rho - x, \pi - \rho)$  for some  $(\rho, \sigma, \tau)$  where the various angles satisfy the conditions of Theorem 1, bisect  $CD$  at  $H$  and  $CF$  at  $J$ , then check (as above for  $\Delta ABC$ ) whether or not  $\Delta HJF \in (\rho, \sigma, \tau)$  satisfies the conditions of Theorem 1 with  $HF \geq FJ \geq JH$ .

We now give a theorem similar to Theorem 1 which deals with the other two similarity classes mentioned in Lemma 1.

*Definition.* The smaller sides ratio  $s$  of a similarity class  $(\rho, \sigma, \tau)$  is obtained by choosing any  $\Delta ABC \in (\rho, \sigma, \tau)$  with  $AB \geq BC \geq AC$ , then setting  $s = BC/AC$ .

**THEOREM 2.** Assume that  $\tau \leq \sigma \leq \rho$ ,  $x + \tau \geq \max\{\sigma, \rho - x\}$ , and  $\pi - \rho \geq \rho - x$ . Then for  $n \geq 1$ , after  $n$  cycles of the bisection method applied to

(i)  $\Delta \in (\rho - x, \sigma, x + \tau)$ , we have  $2^{2n-1}$  triangles in  $(\rho - x, \sigma, x + \tau)$  each with diameter  $d(\Delta)/2^n$  and  $2^{2n-1}$  triangles in  $(x, \tau, \rho + \sigma - x)$  each with diameter  $sd(\Delta)/2^n$ :

(ii)  $\Delta' \in (x, \tau, \rho + \sigma - x)$ , we have  $2^{2n-1}$  triangles in  $(x, \tau, \rho + \sigma - x)$  each with diameter  $d(\Delta')/2^n$  and  $2^{2n-1}$  triangles in  $(\rho - x, \sigma, x + \tau)$  each with diameter  $d(\Delta')/2^n s$ .

Here  $s$  is the smaller sides ratio of  $(\rho, \sigma, \tau)$ .

*Proof.* Analogous to that of Theorem 1.

**COROLLARY 2.** Suppose  $\tau \leq \sigma \leq \rho$ ,  $x + \tau \geq \max\{\sigma, \rho - x\}$ , and  $\pi - \rho \geq \rho - x$ . Then in the notation of the introduction

(i) if  $\Delta_{01} \in (\rho - x, \sigma, x + \tau)$ , then  $m_j \leq s^{(j/2)} d(\Delta_{01})$  for  $j \geq 1$ , with equality for even  $j$ ;

(ii) if  $\Delta_{01} \in (x, \tau, \rho + \sigma - x)$ , then  $m_j \leq (1/2)^{(j/2)} d(\Delta_{01})$  for  $j \geq 1$ , with equality for even  $j$ .

Note that since we are assuming that  $AC \geq \max\{AB/2, CD\}$  in Figure 1, we have  $1 \leq s \leq 2$ .

*Remark.* Given a triangle  $\Delta RST$  with  $RS = d(RST)$ , to decide whether or not  $\Delta RST \in (\rho - x, \sigma, x + \tau)$  or  $(x, \tau, \rho + \sigma - x)$  for some  $(\rho, \sigma, \tau)$  where the various angles satisfy the conditions of Theorem 2, bisect  $RS$  at  $W$  (say). Examine the triangles  $\Delta RWT$  and  $\Delta WST$ . If one of these is in  $(\rho, \sigma, \tau)$ , where the angles satisfy the conditions of Theorem 2, then

(i)  $2WT \geq RS \Rightarrow \Delta RST$  is in the corresponding  $(\rho - x, \sigma, x + \tau)$ ,

(ii)  $2WT \leq RS \Rightarrow \Delta RST$  is in the corresponding  $(x, \tau, \rho + \sigma - x)$ .

Various conditions sufficient for a given triangle  $\Delta XYZ$  to lie in one of the four sets of similarity classes considered can be obtained by elementary calculations using the cosine rule for triangles. For example, given  $\Delta XYZ \in (\alpha, \beta, \gamma)$  with  $XY \geq YZ \geq XZ$  and  $\alpha \geq \beta \geq \gamma$ , then

(i) if  $\cos \gamma \leq 3/4$ , then  $\Delta XYZ \in (\rho, \sigma, \tau)$  satisfies the conditions of Theorem 1;

(ii) if  $XY/YZ \geq 2/\sqrt{3}$  and  $\cos \gamma \leq \sqrt{3}/2$ , then  $\Delta XYZ \in (\rho, \sigma, \tau)$  and  $\Delta XYZ \in (x, \rho - x, \pi - \rho)$ , both satisfying the conditions of Theorem 1;

(iii) if  $3/4 \geq \cos \beta \geq \max\{XZ/XY, XY/4XZ\}$ , then  $\Delta XYZ \in (x, \tau, \rho + \sigma - x)$  satisfies the conditions of Theorem 2.

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